



# A new proof of the degeneracy property of the longest-edge $n$ -section refinement scheme for triangular meshes

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## ABSTRACT

In this note, by using complex variable functions, we present a new simpler proof of the degeneracy property of the longest-edge  $n$ -section of triangles for  $n \geq 4$ . This means that the longest-edge  $n$ -section of triangles for  $n \geq 4$  produces a sequence of triangles with minimum interior angle converging to zero.

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## 1. Introduction

Unstructured mesh generation and adaptive mesh refinement methods for two- and three-dimensional complex domains are very successful tools for the efficient solution of numerical application problems. A major drawback of these methods is that they may produce poorly shaped elements causing the numerical solution to be less accurate and more difficult to compute [1,9]. In the area of adaptive finite element methods, mesh refinement algorithms that maintain the *non-degeneracy* of the elements are certainly desirable. Non-degeneracy means that the minimum angle of the triangles is bounded away from zero when the partition or the refinement is applied.

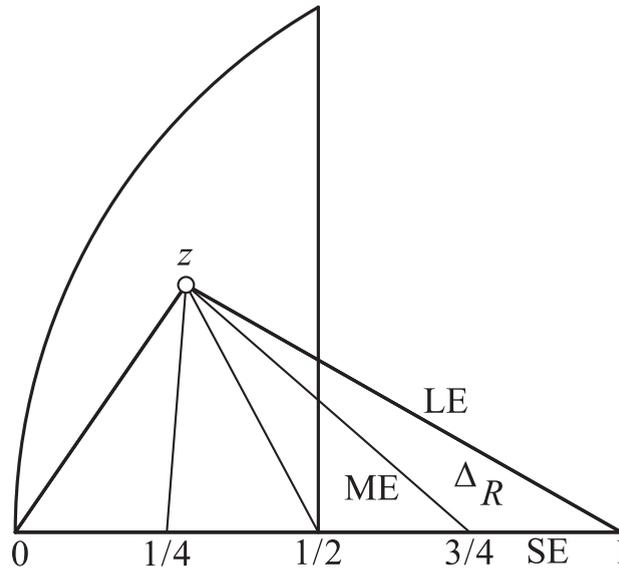
The longest-edge bisection is one possible way of subdividing a triangle and constitutes an efficient scheme for reducing the obtuse angles in mesh refinement [3,6]. The non-degeneracy property of longest-edge bisection of triangles is formulated usually as follows: Let  $t$  be an initial triangle with smallest angle  $\alpha$ . Then the longest-edge partition of triangle  $t$  and its descendants only produces triangles whose smallest angles are always greater or equal to  $\alpha/2$  [7]. For the case of longest-edge trisection of triangles a similar bound has been recently found using a space of triangular shapes endowed with a hyperbolic metric. In fact, for this case the smallest angle produced by longest-edge trisection is always greater or equal to  $\alpha/c$  with  $c = \frac{\pi/3}{\arctan(\sqrt{3}/11)}$  [5,2].

## 2. Longest-edge $n$ -section of the triangles, for $n \geq 4$

In a recent paper [8], Suárez et al. have proved by elementary geometric arguments the degeneracy property of the triangular meshes obtained by iterative application of longest-edge  $n$ -section of the triangles, when  $n \geq 4$ . In this short note, we propose a new simpler proof of the same property by using a normalization process and elementary complex variable functions.

**Theorem 1** [8]. *The iterative application of longest-edge  $n$ -section when  $n \geq 4$  to an initial triangle  $\Delta ABC$  generates a sequence of new triangles in which  $\lim_{k \rightarrow \infty} \alpha_k = 0$ ,  $\alpha_k$  being the minimum triangle angle in iteration  $k$ .*

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**Fig. 1.** Normalized shape region  $\mathfrak{E}$  and longest-edge 4-section of triangle determined by  $z$ : LE, ME and SE are respectively longest, medium and shortest edge of triangle  $\Delta_R$ .

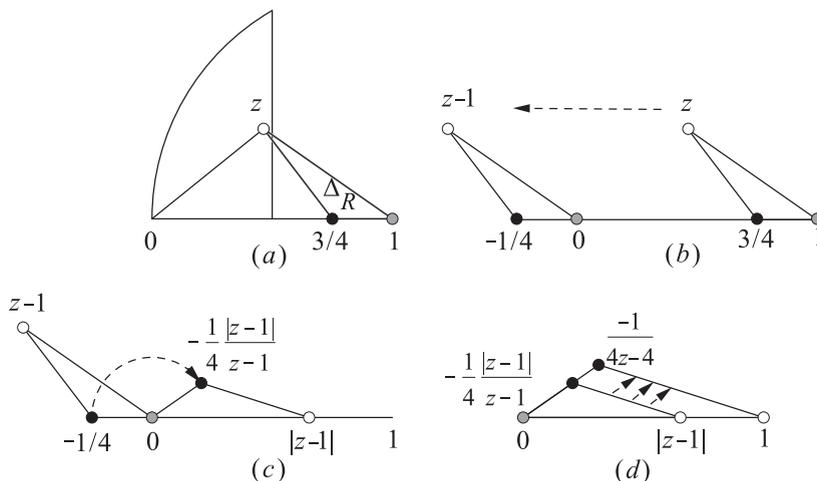
**Proof.** A geometrical diagram to study the shape of triangles generated by iterative application of triangle subdivision has been introduced in [4]. The map diagram provides a convenient way to visualize the evolution and migration of element shapes. We consider all the triangles normalized in such a way that for a given triangle the longest-edge is scaled to have unit length, with extreme points  $(0,0)$  and  $(1,0)$ , and the shortest angle at point  $(1,0)$ . It is clear, that after some Euclidean movements any triangle may be transformed in a similar triangle in that way.

Note that a normalized triangle is determined by the opposite point to the longest-edge. This point  $z$  will be called apex of the triangle and may be considered as a complex point. The region of the plane for the apexes of normalized triangles will be called *normalized shape region* and denoted by  $\mathfrak{E}$ :

$$\mathfrak{E} = \{z \in \mathbb{C}, \text{Im}(z) > 0, \text{Re}(z) \leq 1/2, |z - 1| \leq 1\}.$$

In this form, all similar triangles are represented once normalized by a unique point  $z$  in the normalized shape region. Note that then the shortest angle  $\alpha$  is written as function of apex point  $z$  as  $\alpha = \text{Arg}(1 - \bar{z})$ .

Fig. 1 shows the normalized shape region  $\mathfrak{E}$ , and a triangle divided by the longest-edge 4-section. The sub-triangle with vertex at point  $(1,0)$  is denoted by  $\Delta_R$  in Fig. 1. The shortest edge of  $\Delta_R$ , SE in the figure, has extreme points  $(3/4,0)$  and  $(1,0)$ . In general for the case of longest-edge  $n$ -section, being  $n \geq 4$ , the shortest edge of the sub-triangle with vertex point  $(1,0)$  would be  $1/n$  length.



**Fig. 2.** Normalization of Triangle  $\Delta_R$  for LE 4-section, expressed by compositions of complex variable functions.

Note that triangle  $\Delta_R$  depends on the initial triangle, and so on point  $z$ . Therefore the normalization process of triangle  $\Delta_R$  may be understood as a complex variable function  $w(z)$  being  $z$  the apex point of the initial triangle and  $w(z)$  the apex point of the triangle  $\Delta_R$  once normalized. For the case  $n = 4$ , function  $w(z)$  is obtained in Fig. 2, where  $w(z) = \frac{-1}{4(z-1)}$ . For the general case,  $n \geq 4$ , it would be  $w(z) = \frac{-1}{n(z-1)}$ .

Note that since  $|z - 1| \geq \frac{1}{2}$ , then  $\frac{\text{Im}w(z)}{\text{Im}z} \leq \frac{2}{n}$ .

$$\frac{\text{Im}w(z)}{\text{Im}z} = \frac{\text{Im} \frac{-1}{n(z-1)}}{\text{Im}z} = \frac{\text{Im} \frac{-(\bar{z}-1)}{n(z-1)}}{\text{Im}z} \leq \frac{2}{n} \frac{\text{Im}(1 - \bar{z})}{\text{Im}z} = \frac{2}{n}.$$

Therefore, after  $k$  compositions of functions  $w$  we get  $\text{Im}w^k(z) \leq \left(\frac{2}{n}\right)^k \text{Im}z$ . In other words,  $\text{Im}w^k(z)$  tends to zero as  $k$  goes to infinity, and so the argument of  $1 - w^k(z)$ . This completes the proof.  $\square$

Note that, since the estimate  $\frac{\text{Im}w(z)}{\text{Im}z} \leq \frac{2}{n}$  does not depend on point  $z$ , it also holds, for example, if  $z$  belongs to the boundary of the normalized region  $\mathfrak{E}$ . For instance, let  $z = \frac{1}{2} + hi$  then the minimum initial angle is  $\alpha = \arctan 2h$  while the minimum angle of  $w(z)$  is  $\alpha_{w(z)} = \arctan \frac{h}{\frac{n-2}{4} + nh^2}$  which tends to zero as  $h$  tends to zero.

### 3. Final remarks

Here the degeneracy property of the triangular partitions obtained by longest-edge  $n$ -section, for  $n \geq 4$ , of triangles has been proved. It has been proved that if  $\alpha_k$  is the minimum angle of the triangles obtained by the LE- $n$ -section of an initial triangle  $t$ , then  $\alpha_k$  tends to zero for  $k \rightarrow \infty$ . This fact is in contrast with the non-degeneracy property of the longest-edge  $n$ -section for  $n = 2, 3$ .

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